

Introduction to linear programming

Lycée Jules Haag, Besançon, April 12th 2012

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How many boxes of each kind should they make in order to maximize the profit ?

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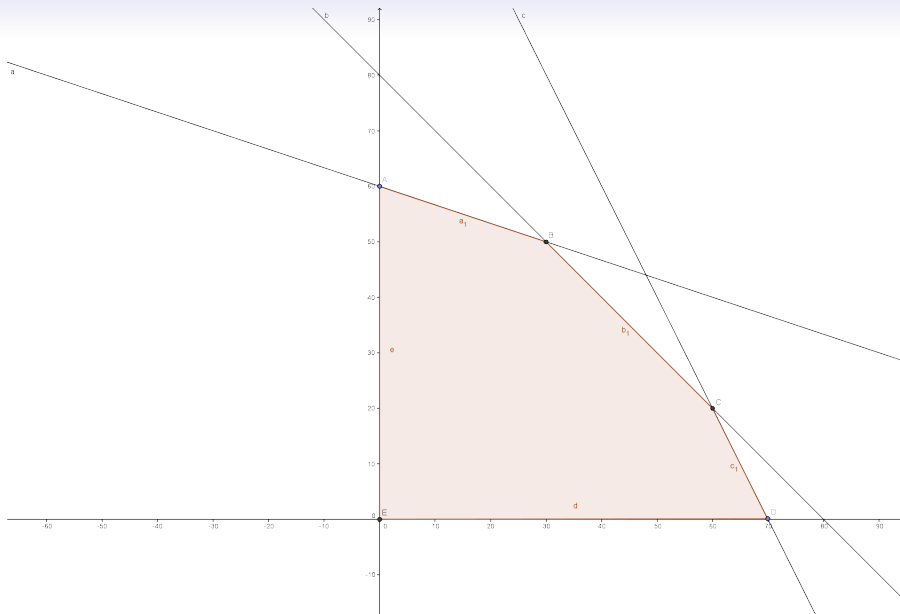
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This will be called the **feasible region**.



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Can we **conjecture** the statement of the general theorem ?

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Step 1 The maximum of the objective function on a **segment** is attained at an **endpoint** of the segment.

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We obtain

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This yields the conclusion of Step 1.

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End of proof.

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Why should we prove it ???